

# Delay-Optimal Cross-Layer Design for Wireless Systems

**Vincent Lau**

Dept of ECE

Hong Kong University of Science and Technology

May 2009

# Outline

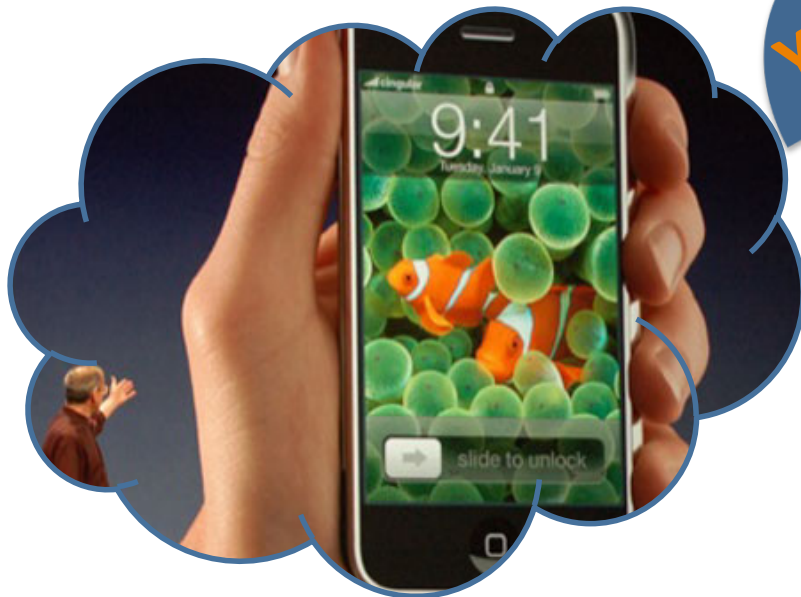
---

- 💣 Introduction and Motivation
- 💣 Survey of Existing Approaches
- 💣 Example:
  - Distributive Delay-Optimal Control for Uplink OFDMA via
  - Localized Stochastic Learning and Auction Game
- 💣 Convergence Analysis
- 💣 Asymptotic Optimality
- 💣 Conclusion

# Introduction and Motivation

## 💣 Why delay performance is important?

- 💣 “WHAT??!! He is stuck in the air!! !\$\*(&#%\*!(!”
- 💣 “You must be kidding me! Buffering at such an important moment!!??”



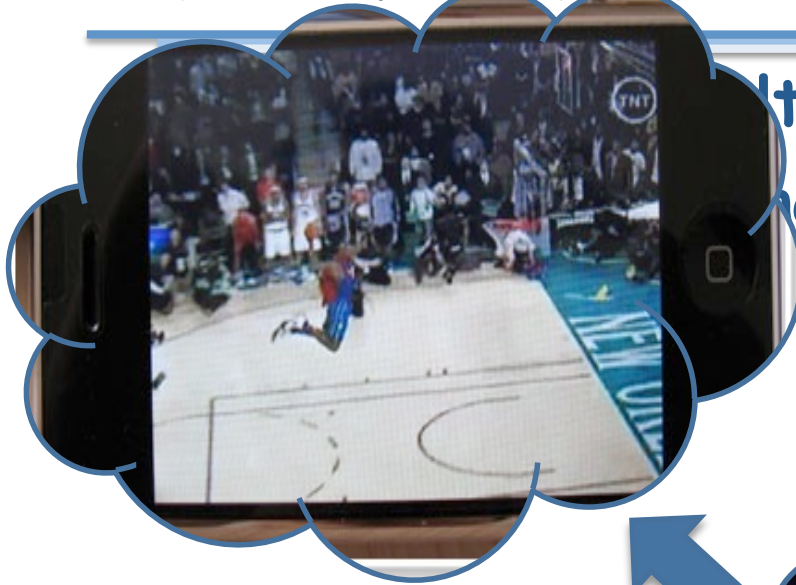
You Tube



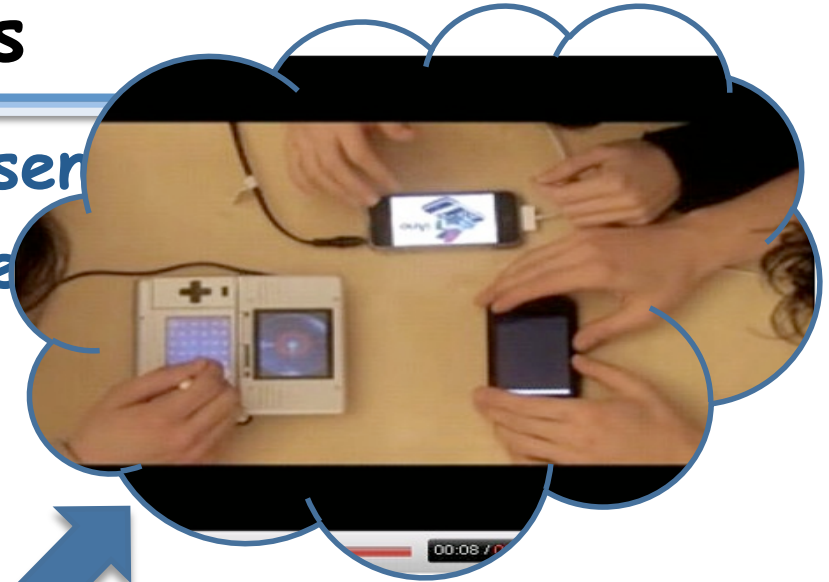
**Fact I:**  
Real-life applications are **delay-sensitive**

# Introduction and Motivations

Multiple delay-sensitive applications at different times



Keep track of a game



Play multi-player game



**Fact II:**  
Different users and applications have **heterogeneous delay requirements**

Keep talking to some friends





# Related Works

## 💣 OFDMA Joint Power and Subband Design for PHY Performance

[Yu'02], [Hoo'04],[Seong'06], etc.

- Selects the strongest user per subband
- Time-Frequency Water-filling Power Allocation
- Assuming knowledge of **perfect CSIT**.

[Lau'05], [Wong'09], [Brah'07] etc.

- Robust Power and Subband Control with **limited feedback** or **outdated CSIT (packet errors)**.

### Remark:

Only adapt based on **CSIT**, ignoring queue states and optimize **PHY layer performance only (throughput or PFS)**

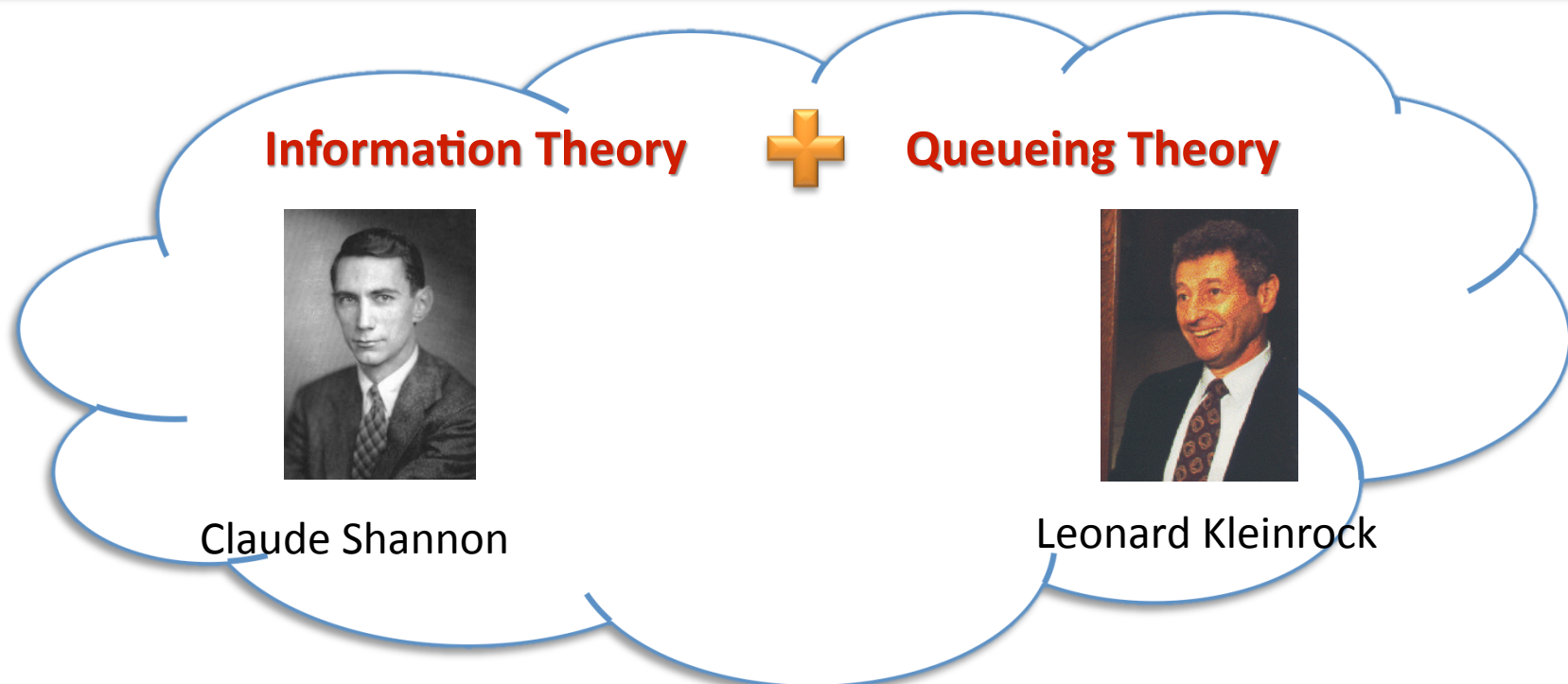
**Conclusion:** Very important to make use of both **(channel state info) CSI** and **(queue state info) QSI** for delay sensitive applications

# Introduction and Motivations

## 💣 Challenges to incorporate QSI and CSI in adaptation

**Challenge 1:** Requires both **Information theory** (modeling of the PHY dynamics) & **Queueing theory** (modeling of the delay/buffer dynamics)

**Challenge 2:** Brute-force approach cannot lead to any viable solution



When Shannon meets Kleinrock...

# Existing Approaches to deal with Delay-Optimal Control

## 💣 Various approaches dealing with delay problems

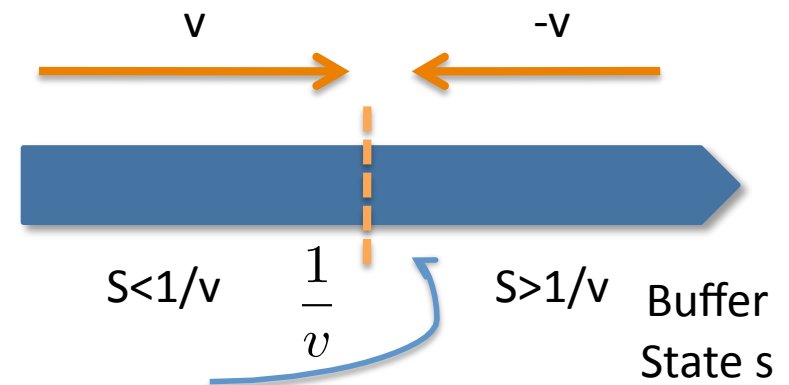
### Approach I : Stability Region and Lyapunov Drift [Berry'02], [Neely'07], etc.

- Discuss **stability region** of point-to-point SISO and multiuser SISO.
- Also considered asymptotically delay-optimal control policy based on “**Lyapunov Drift**”
- The authors obtained interesting tradeoff results as well as insight into the structure of the **optimal control policy at large delay regime**.

#### Remark:

This approach allows simple control policy with design insights but the control will be good only for asymptotically **large delay regime**.

#### Buffer Partitioning



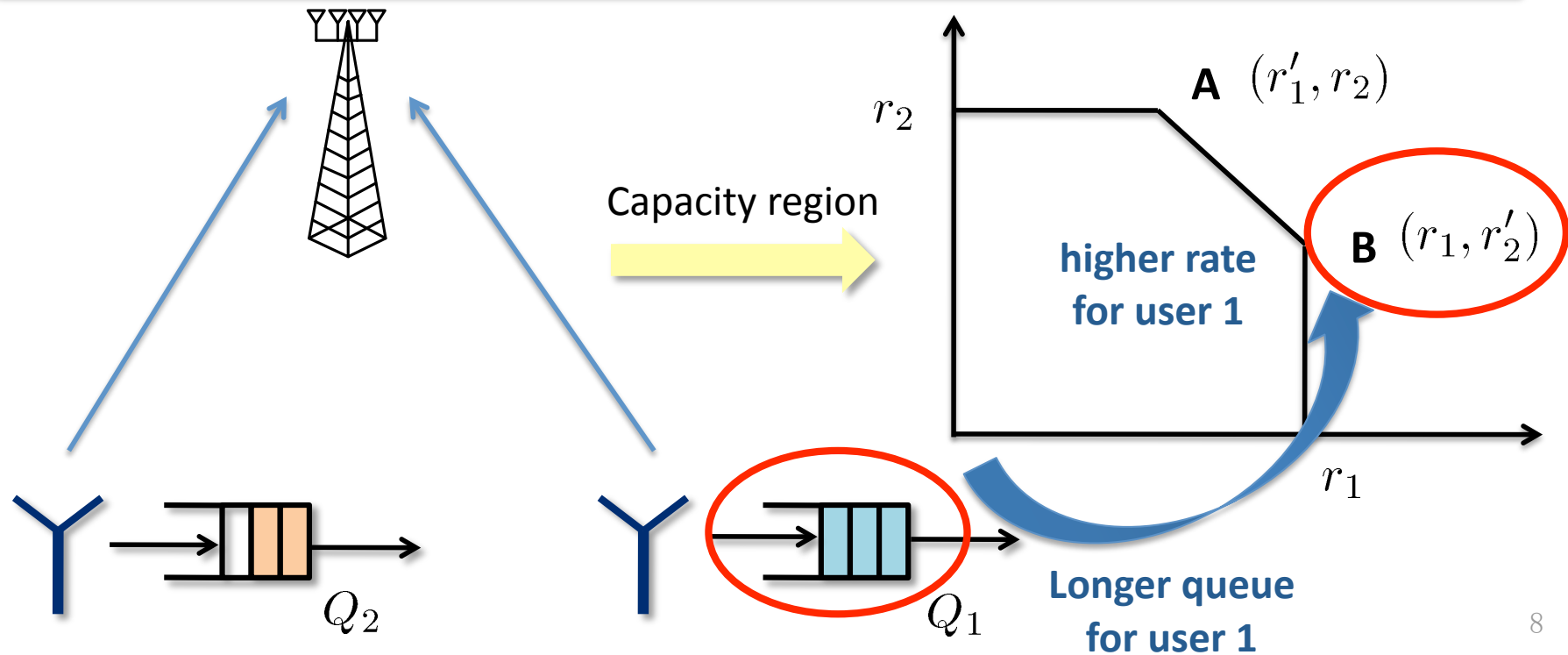
To regulate the buffer state towards  $1/v$

# Related Works

## 💣 Various approaches dealing with delay problems

### Approach II [Yeh'01PhD], [Yeh'03ISIT]

- Symmetric and homogeneous users in multi-access fading channels
- Using **stochastic majorization theory**, the authors showed that the **longest queue highest possible rate (LQHPR)** policy is delay-optimal



# Related Works

## 💣 Various approaches dealing with delay problems

**Approach III :** [Hui'07], [Tang'07], etc.

To **convert the delay constraint into average rate constraint** using tail probability at large delay regime (large deviation theory) and solve the optimization problem using information theoretical formulation based on the rate constraint.

### **Remark:**

While this approach allows potentially simple solution, the control policy will be **a function of CSIT only** and such control will be good only for **large delay regime**.

### **Note:**

In general, the delay-optimal power and precoder adaptation should be a function of **both the CSI and the QSI**.

# Related Works

## 💣 Various approaches dealing with delay problems

### Approach IV : [Bertsekas'87]

The problem of finding the optimal control policy (to minimize delay) is casted into a **Markov Decision Problem (MDP)** or a stochastic control problem.

### Remark:

- Unfortunately, it is well-known that there is no easy solution to MDPs in general.
- Brute-force **value iteration** and **policy iteration** are **very complex and time-consuming**.
- The curse of **dimensionality!!**



# Technical Challenges To be Solved

---

## Challenge 1:

A systematic approach for **low complexity** delay-optimal control policy in **general delay regime**.

## Challenge 2:

Curse of Dimensionality: Exponential Complexity due to coupling among multiple delay-sensitive **heterogeneous users**.

## Challenge 3:

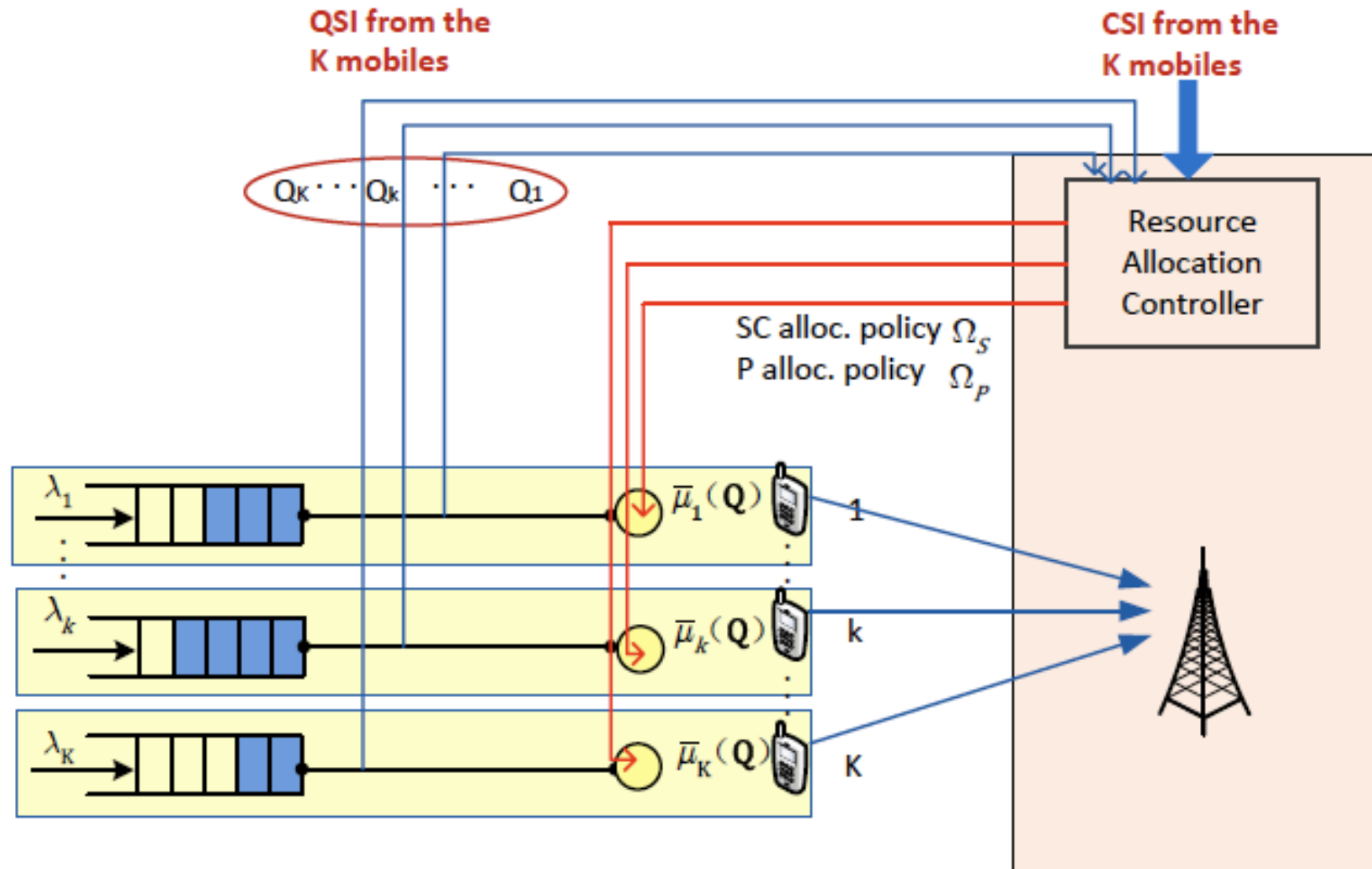
Structure of the delay-optimal policy.

## Challenge 4:

Distributive Implementation (Function of Local CSI and Local QSI, e.g. uplink OFDMA)

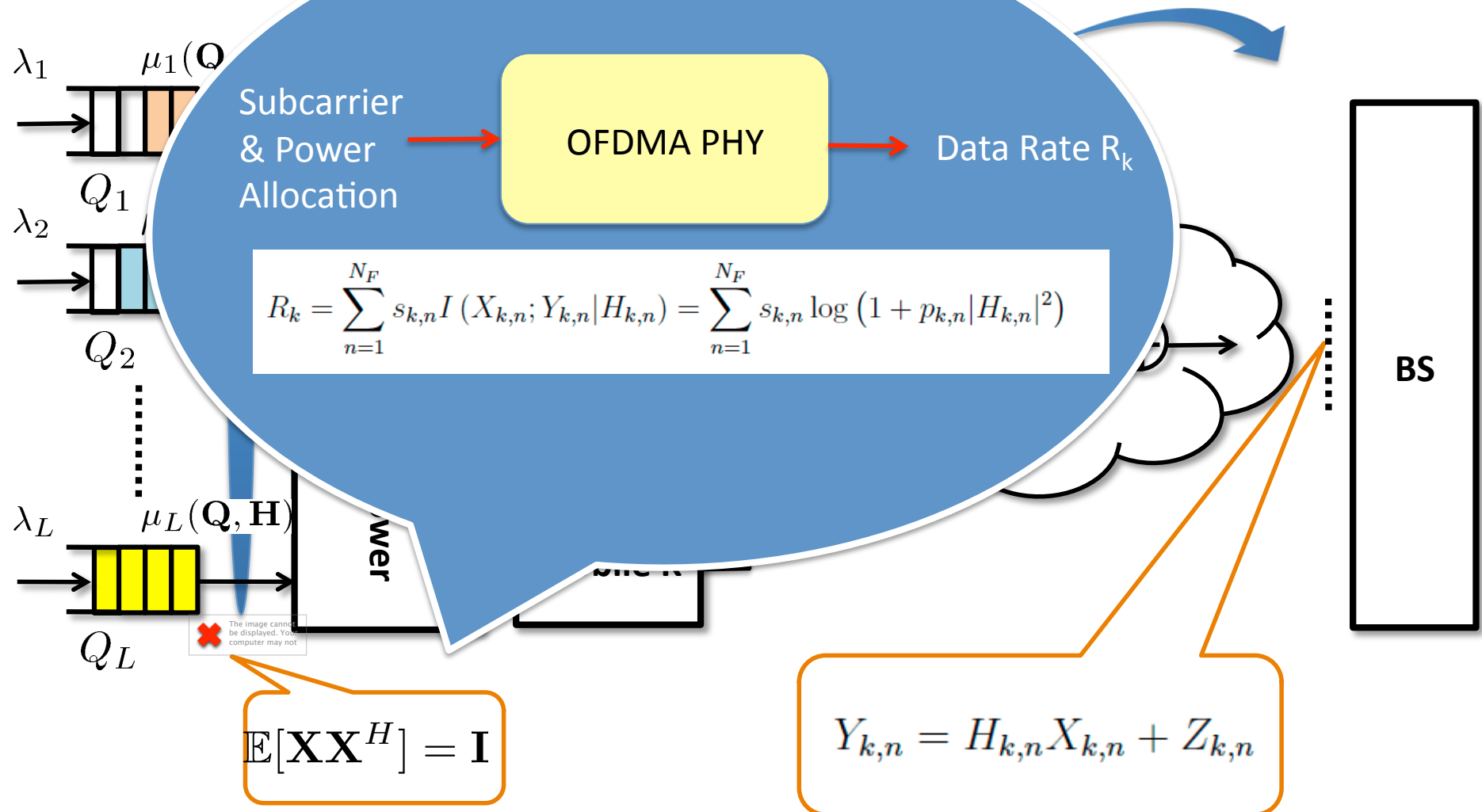
Example :  
Distributive Delay-Optimal Control  
for Uplink OFDMA via Localized  
Stochastic Learning and Auction Game

# Uplink OFDMA System Model

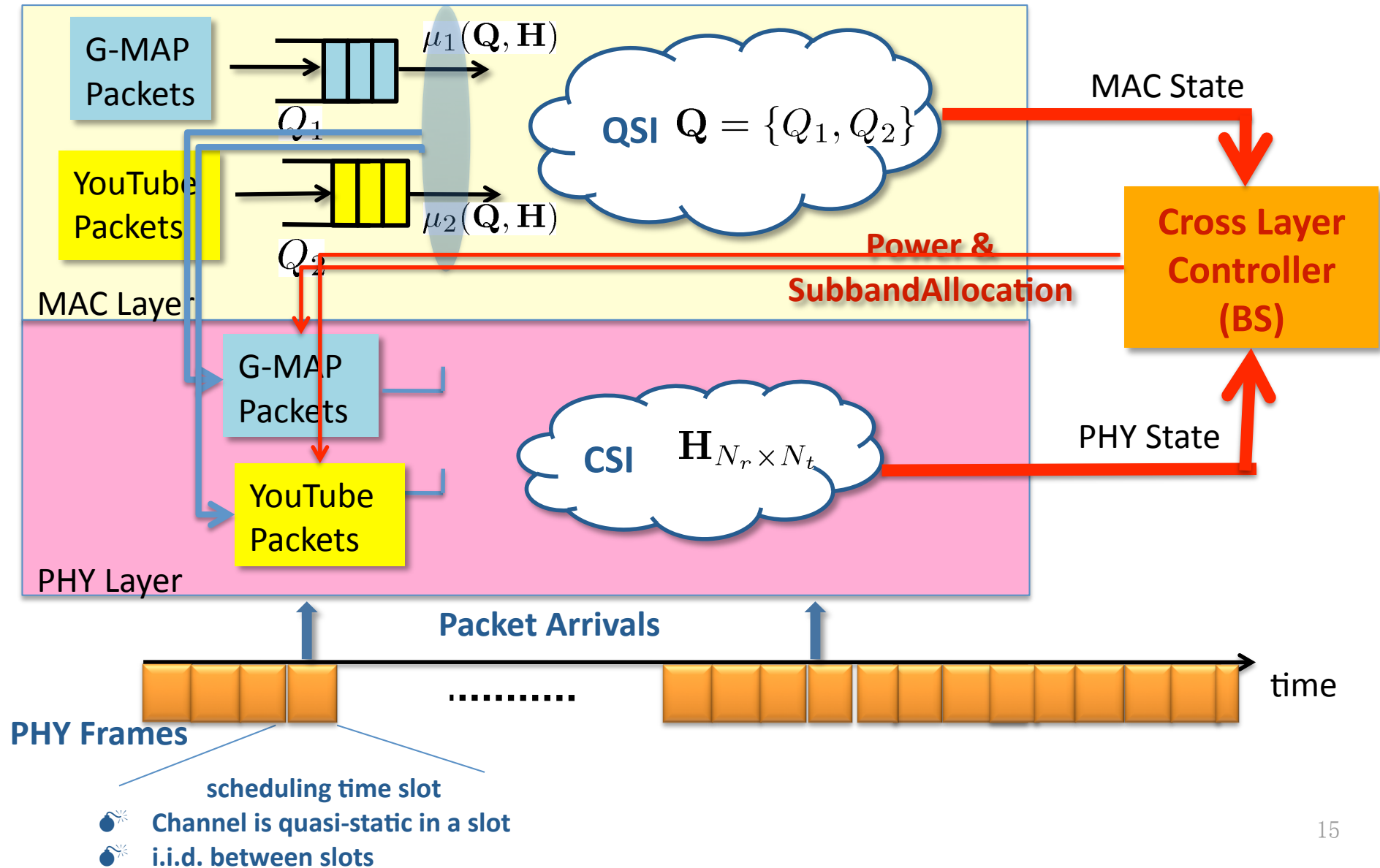


# OFDMA PHY Model

## OFDMA Physical Layer



# Source Model and System States



# OFDMA Queue Dynamics

- Time domain partitioned into scheduling slots
- CSI  $H(t)$  remains quasi-static within a slot and is i.i.d. between slots
- Packet arrival  $A(t)=(A_1(t), \dots, A_K(t))$  where  $A_k(t)$  i.i.d. according to a general distribution  $P(A)$ .
- $N_k(t)$  denotes the random packet size, i.i.d.
- $Q_k(t)$  denotes the number of packets waiting in the  $k$ -th buffer at the  $t$ -th slot.

$$Q_k(t+1) = \min\{[Q_k(t) - R_k(t)\tau / N_k(t)]^+ + A_k(t), N_Q\}$$

- Global System State (CSI, QSI)**  $\chi(t) = (H(t), Q(t))$ 
  - Total number of bits Transmitted in the  $t$ -th slot



# OFDMA Delay-Optimal Formulation

## Stationary Power and Subband Allocation Control Policy

- 💣 A mapping  $\Omega = (\Omega_p, \Omega_s)$  from the system state  $\chi$  to a power and subband allocation actions.

$$\Omega_p(\chi) = \{p_{k,n}\}, \quad \Omega_s(\chi) = \{s_{k,n}\}$$

$$\sum_{n=1}^{N_F} E[p_{k,n}] \leq P_k \quad \forall k \in \{1, K\}, \quad p_{k,n} \geq 0 \quad \text{(Power Constraint)}$$

$$\sum_{k=1}^K s_{k,n} = 1 \quad \forall n \in \{1, N_F\} \quad \text{(Subband Allocation Constraint)}$$

$$\Pr[Q_k = N_Q] \leq P_k^d \quad \forall k \in \{1, K\} \quad \text{(Packet Drop Rate Constraint)}$$

# OFDMA Delay-Optimal Formulation

Definitions: Average Delay, Power and Packet Drop Constraints under a control policy  $\Omega$

$$\bar{T}_k(\Omega) = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \mathbb{E}[Q_k(t)] = \mathbb{E}_{\pi_\chi} [Q_k] \quad \forall k \in \{1, K\}$$

$$\bar{P}_k(\Omega) = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \mathbb{E} \left[ \sum_n p_{k,n}(t) \right] = \mathbb{E}_{\pi_\chi} \left[ \sum_n p_{k,n} \right] \leq P_k \quad \forall k \in \{1, K\}$$

$$\bar{P}_k^d(\Omega) = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \mathbb{E} \left[ \mathbf{1}[Q_k(t) = N_Q] \right] = \mathbb{E}_{\pi_\chi} \left[ \mathbf{1}[Q_k = N_Q] \right] \leq P_k^d \quad \forall k \in \{1, K\}$$

$\mathbb{E}_{\pi_\chi}$  denotes expectation w.r.t. the underlying measure  $\pi$ .

Little's Law: average no. of packets = average arrival rate \* average delay  
*the average delay (in terms of seconds)  $\propto$  the average queue length*

# Formulation

“Positive Weighting Factor”

$$\beta = (\beta_1, \dots, \beta_K)$$

Pareto Optimal delay boundary

$$d(\chi, \{p, s\}) = \sum_k \beta_k Q_k$$

$$\min_{\Omega} J_{\beta}^{\Omega} \equiv \sum_{k=1}^K \beta_k \bar{T}_k(\Omega) = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \mathbb{E} \left[ d(\chi(t), \Omega(\chi(t))) \right]$$

subject to the power and packet drop rate constraints

**Solution: Markov Decision Problem (MDP)**

**Key Idea: Divide-and-Conquer**

To break a large problem (optimization over the whole policy space) into smaller sub problems (optimization over a control action at a stage).

# Overview of Markov Decision Problem Formulation

## 💣 Specification of an Infinite Horizon Markov Decision Problem

- Decisions are made at points of time – decision epochs
- **System state and Control Action Space:**
  - At the  $t$ -th decision epoch, the system occupies a state  $S_t$
  - The controller observes the current state and applies an action  $A_t$
- **Per-stage Reward & Transition Probability**
  - By choosing action  $A_t$  the system receives a reward  $R(S_t, A_t)$
  - The system state at the next epoch is determined by a transition probability kernel  $\Pr(S_{t+1} \mid S_t, A_t)$
- **Stationary Control Policy:**
  - The set of actions for all system state realizations  $A_t = \pi(S_t)$
- **The Optimization Problem:**
  - Average Reward
  - Optimal Policy

$$\bar{R}^* = \max_{\pi} \lim_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[ \sum_{t=1}^T R(S_t, A_t) \right]$$

# Overview of Markov Decision Problem Formulation

## 💣 Solution of an Markov Decision Problem

### Key Criterion: Bellman's Equation

Under some technical conditions, the optimal value of the problem is given by the solution of the **Bellman's Equation**.

## 💣 Optimal average reward

$$\overline{R}^* = \theta$$

## 💣 Optimal policy (Fixed Point Problem on Functional Space)

$$\pi^* = \arg \max_{A_i} \left\{ r(S_i, A_i) + \sum_{S_j} Pr(S_j | S_i, A_i) V(S_j) \right\}$$

# Constrained Markov Decision Problem Formulation

$$d(\chi, \{p, s\}) = \sum_k \beta_k Q_k$$

**CMDP Formulation:** Find the optimal control policy  $\Omega$  that minimizes

$$\min_{\Omega} J_{\beta}^{\Omega} = \sum_{k=1}^K \beta_k \bar{T}_k(\Omega) = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \mathbb{E} \left[ d(\chi(t), \Omega(\chi(t))) \right]$$

subject to the power and packet drop rate constraints

**Lagrangian approach to the Constrained MDP:**

$$\min_{\Omega} L_{\beta}^{\Omega}(\gamma) = \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \mathbb{E} \left[ g(\chi(t), \Omega(\chi(t))) \right]$$

$$g(\gamma, \chi, \Omega(\chi)) = \sum_k \left( \beta_k Q_k + \bar{\gamma}^k (\sum_n p_{k,n} - P_k) + \underline{\gamma}^k (1[Q_k = N_Q] - P_k^d) \right)$$



# Optimal Solution

## 💣 Infinite Horizon Average Reward MDP

- 💣 Given a stationary control policy  $\Omega$ ,  
 $\{\chi(t), g(\chi(t), \Omega(\chi(t)))\}$  evolves like a Markov Chain  
with transition kernel:

$$\Pr[\chi(t+1)|\chi(t), \Omega(\chi(t))] = \Pr[\mathbf{H}(t+1)] \Pr[\mathbf{Q}(t+1)|\chi(t), \Omega(\chi(t))]$$

- 💣 Solution is given by the “Bellman Equation”

$$\theta + V(\chi^i) = \min_{u(\chi^i)} \left[ g(\chi^i, u(\chi^i)) + \sum_{\chi^j} \Pr[\chi^j|\chi^i, u(\chi^i)] V(\chi^j) \right]$$

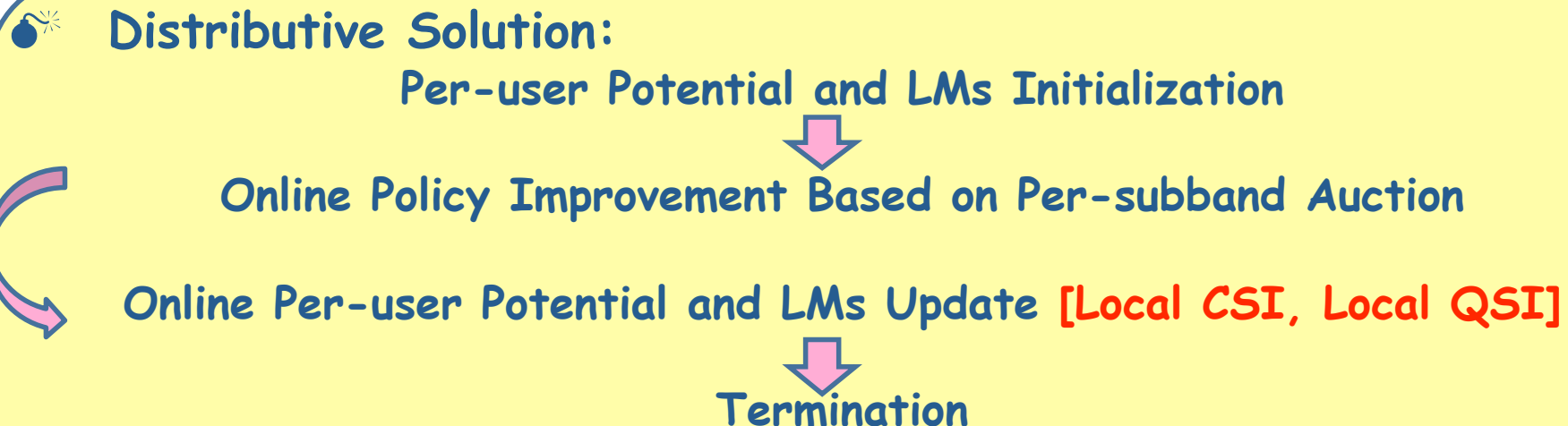
“Potential function” (contribution of the state  $i$  to the average reward)

“Optimal Value”  $\theta = J_{\beta}^* = \inf_{\Omega} J_{\beta}^{\Omega}$

$(N_Q + 1)^K$  Equations and  
unknowns

# Optimal Solution - Online Learning

- 💣 How to determine the potential function ?
  - 💣 Brute-Force solution of the Bellman Equation ? (Value Iteration):
    - 💣 Too complicated, exponential complexity and memory requirement
  - 💣 Online stochastic learning !
    - 💣 Iteratively estimate potential function based on real time observation of CSI and QSI - online value iteration



# Decentralized Solution (I)

## Online Per-user Primal-Dual Potential Learning Algorithm via Stochastic Approximation

$$\hat{V}_{l+1}^k(Q_k^i) = \begin{cases} \hat{V}_l^k(Q_k^i) + \epsilon_l^v \left[ g_k(\gamma_l^k, \chi_k(l)) + \hat{V}_l^k(Q_k(l+1)) \right. \\ \quad \left. - (g_k(\gamma_l^k, \chi_k^I) + \hat{V}_l^k(\bar{Q}_k^I) - \hat{V}_l^k(Q_k^I)) - \hat{V}_l^k(Q_k^i) \right] & \text{if } Q_k^i = Q_k(l) \\ \hat{V}_l^k(Q_k^i) & \text{if } Q_k^i \neq Q_k(l) \end{cases}$$

$$\bar{\gamma}_{l+1}^k = \Gamma(\bar{\gamma}_l^k + \epsilon_l^\gamma (\sum_{n=1}^N p_{k,n}^l - P_k))$$

$$\underline{\gamma}_{l+1}^k = \Gamma(\underline{\gamma}_l^k + \epsilon_l^\gamma (\mathbf{1}[Q_k(l) = N_Q] - P_k^d))$$

**New Observation at the beginning of the (l+1)-th slot**

**Remark (Co**  
**Deterministic N**  
CSI coherence tim

$$\sum_l \epsilon_l^v = \infty, \epsilon_l^v \geq 0, \epsilon_l^v \rightarrow 0, \sum_l \epsilon_l^\gamma = \infty, \epsilon_l^\gamma \geq 0, \epsilon_l^\gamma \rightarrow 0$$

$$\sum_l ((\epsilon_l^v)^2 + 2(\epsilon_l^\gamma)^2) < \infty, \frac{\epsilon_l^v}{\epsilon_l^\gamma} \rightarrow 0$$

Both the per-user potential and 2 LMs are updated simultaneously.

es evolves in the same better solution (no

# Decentralized Solution (II)

- Per-stage auction with K bidders (MSs) and one auctioneer (BS)

- Low complexity Scalarized Per-Subband Auction

- Bidding: Each user submits a bid  $\hat{X}_{k,n}$

- Subband allocation:  $s_{k,n}(\mathbf{H}_n, \mathbf{Q}^i) = \begin{cases} 1, & \text{if } k = k_n^* \text{ and } \hat{X}_{k_n^*,n} > 0 \\ 0, & \text{otherwise } k_n^* = \arg \max_k \hat{X}_{k,n} \end{cases}$

- Power allocation:  $p_{k,n}(\mathbf{H}_n, \mathbf{Q}^i) = s_{k,n}(\mathbf{H}_n, \mathbf{Q}^i) \left( \frac{\bar{N}_k^{OV}(\mathbf{Q}_k)}{\gamma^k} - \frac{1}{|H_{k,n}|^2} \right)^+$

- Charging:  $C_{k,n} = s_{\bar{k}_n^*,n} \hat{X}_{\bar{k}_n^*,n} 1\{k = \bar{k}_n^*\} \quad \bar{k}_n^* = \arg \max_{k \neq k_n^*} \hat{X}_{k,n}$

- Lemma: The per-stage serial optimal scalarized bid**

Water-level depends on QSI (via potential function)

$$\delta \tilde{V}^k(Q_k^i) = \tilde{V}(Q_1^i, \dots, Q_k^i, \dots, Q_K^i) - \tilde{V}(Q_1^i, \dots, [Q_k^i - 1]^+, \dots, Q_K^i)$$

$\hat{X}_{k,n}$

$$\left( \frac{1}{|H_{k,n}|^2} \right)^+$$

# Decentralized Solution

Theorem (**Convergence of online per-user learning**) Under some mild conditions, the distributive learning converges

Remark (Comparison to conventional stochastic learning)

**Conventional SL:** (1) for unconstrained MDP only or LM for CMDP are determined offline by simulation; (2) designed for centralized solution with control action determined entirely from the potential update → Convergence Proof based on standard “**contraction Mapping**” and **Fixed-Point Theorem** argument.

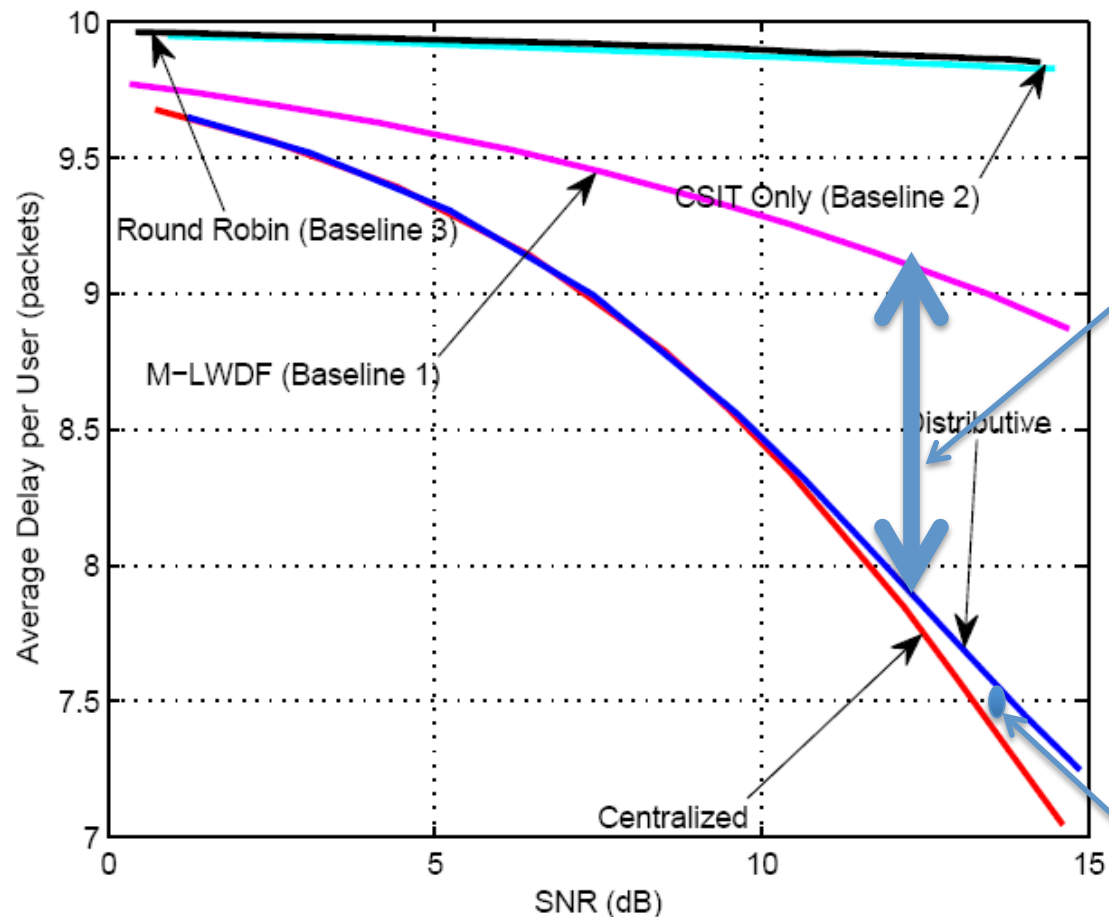
**Proposed SL:** (1) simultaneous update of LM and the potential function; (2) control action is determined by all the users’ potential via per-stage auction → per-user potential update is **NOT a contraction mapping** & standard proof does not apply.

**Bellman equation.**

# Numerical Results

## Average Delay per user vs SNR

The number of users  $K = 2$ , the buffer size  $N_Q = 10$ , the mean packet size  $\bar{N}_k = 305.2$  Kbyte/pck, the average arrival rate  $\lambda_k = 20$  pck/s.



Huge gain in delay performance compared with Modified-Largest Weighted Delay First (M-LWDF), which is the queue length weighted throughput maximization.

Close-to-optimal performance even for small number of users<sup>28</sup>

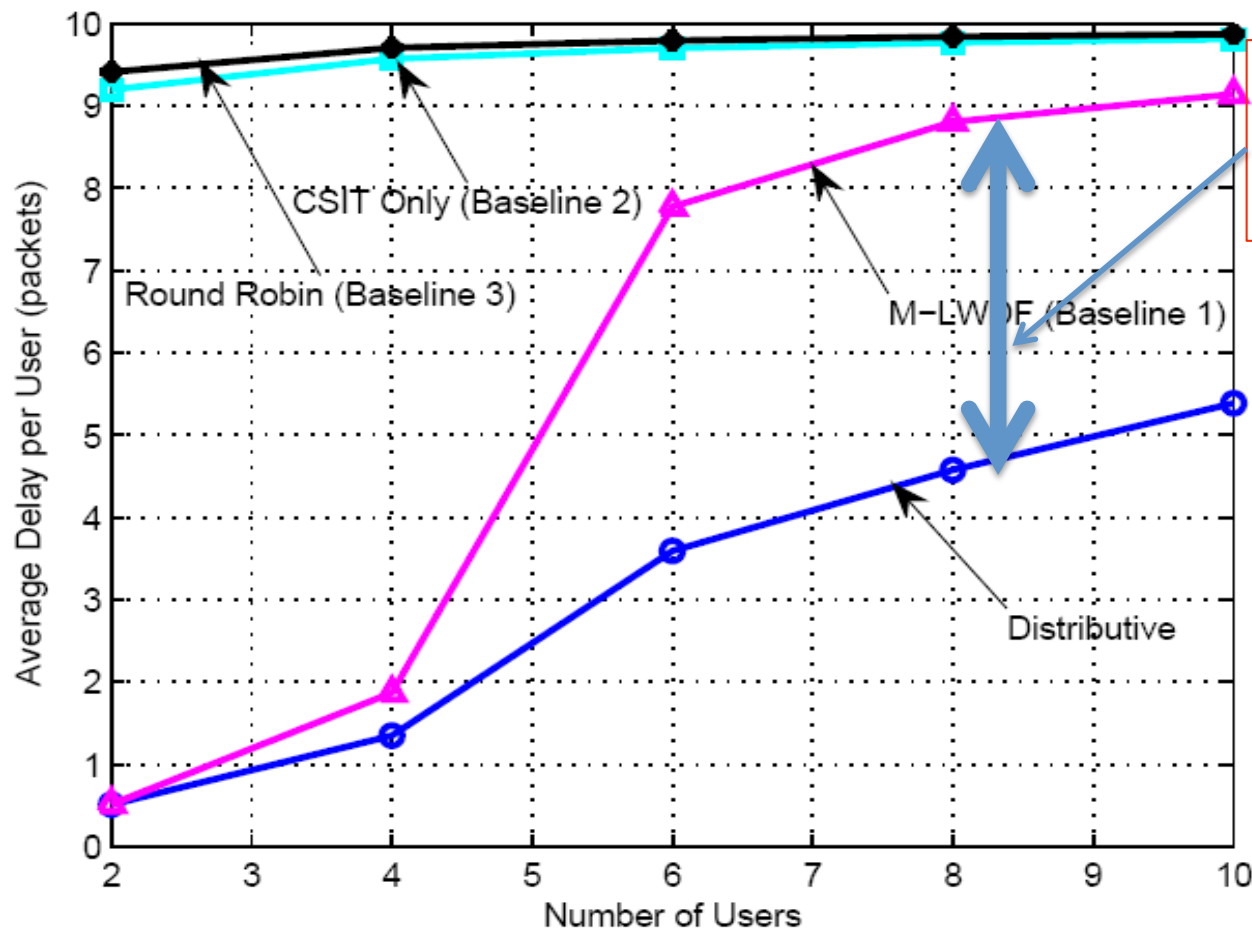


# Numerical Results

## Average Delay per user vs No. of users

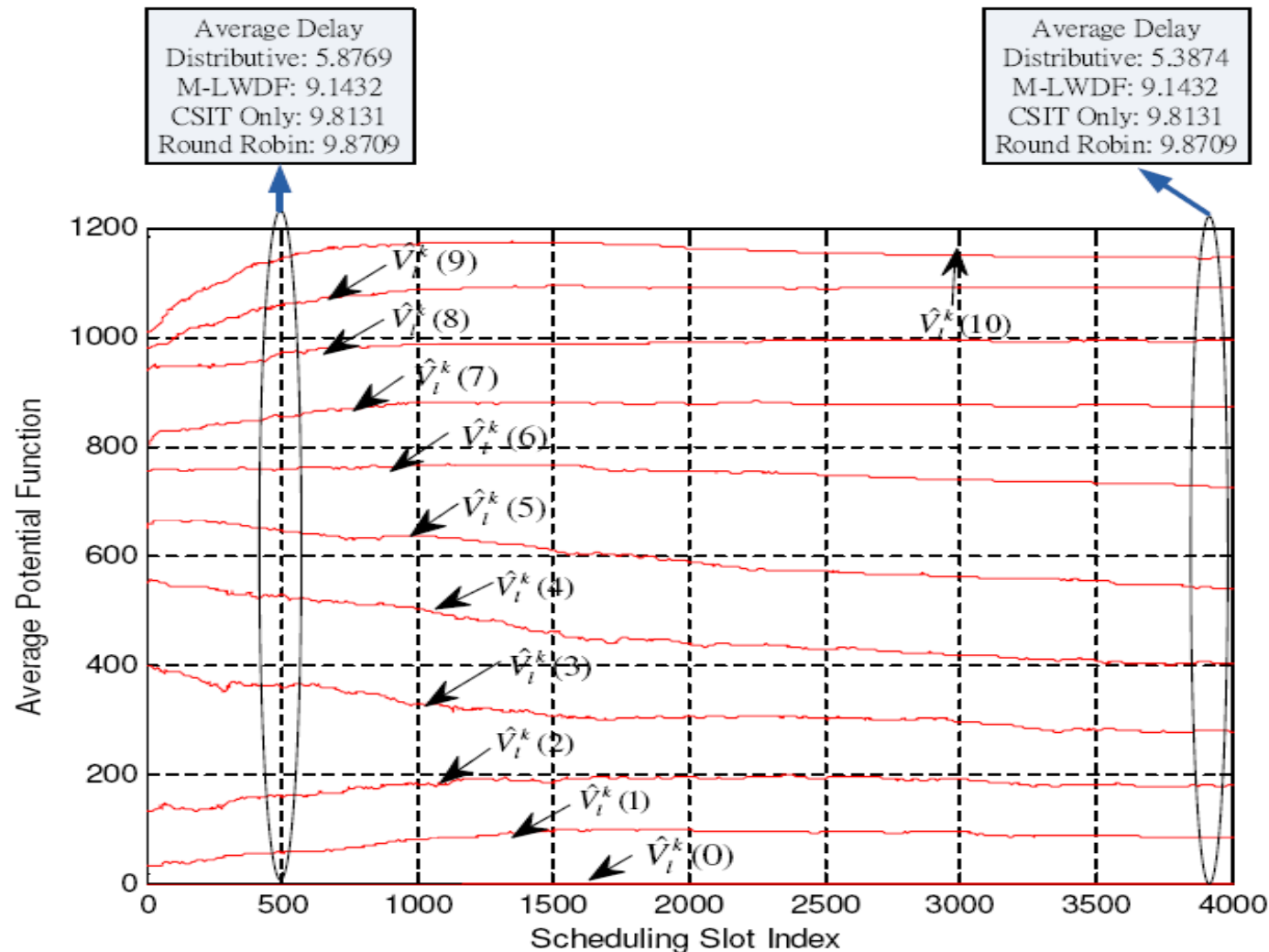
The buffer size  $N_Q = 10$ , the mean packet size  $\bar{N}_k = 78.125$  Kbyte/pck

the average arrival rate  $\lambda_k = 20$  pck/s, the queue weight  $\beta_k = 1$  at a transmit SNR= 10dB.



# Numerical Results

Illustration of convergence property:  
Potential function vs. the scheduling slot index (K=10)



# Conclusion

Distributive Implementation via Decentralized Stochastic Learning and Auction Game

Online Per-user Learning:  
Simultaneous update of LMs and Potentials.  
Almost sure convergence

Optimal Strategy for the Auction Game:  
Delay-Optimal Power Control: Multi-Level Water-Filling  
(QSI  $\rightarrow$  water level; CSI  $\rightarrow$  instantaneous allocation)  
Delay-Optimal Subband Allocation: User selection based on (QSI, CSI)

Asymptotically Global Optimal for large K

# References

---

- V.K.N.Lau,Y.Chen, “***Delay-Optimal Precoder Design for Multi-Stream MIMO System***”, to appear *IEEE Transactions on Wireless Communications*, May 2009.
- V.K.N.Lau, Y. Cui, “***Delay Optimal Power and Subcarrier Allocation for OFDMA System via Stochastic Approximation***”, submitted to *IEEE Transactions on Wireless Communication*, 2008.
- K.B. Huang, V.K.N.Lau, “***Stability and Delay of Zero-Forcing SDMA with Limited Feedback***”, submitted to *IEEE Transactions on Information Theory*, Feb. 2009.
- L.Z. Ruan, V.K.N.Lau, “***Multi-level Water-Filling Power Control for Delay-Optimal SDMA Systems***”, submitted to *IEEE Transactions on Wireless Communication*, 2008.

Thank you!

Questions are Welcomed!

Vincent Lau - [eeKNLau@ee.ust.hk](mailto:eeKNLau@ee.ust.hk)

<http://www.ee.ust.hk/~eeKNLau>